

On the clonal approach in the mathematical theory of social choice

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Arrow's impossibility theorem

Arrow's impossibility theorem. Notation

- A – a non-empty (finite) set (of alternatives);
- n – a natural number (of voters), $n \geq 1$;
- (*individual preferences*) = (strict) linear order on A ;
- $\text{Ord}(A)$ – the set of all (strict) linear orders on A ;
- *profile* = n -tuple of linear orders on A ;
- (*universal aggregation rule*) = function $f : (\text{Ord}(A))^n \rightarrow \text{Ord}(A)$.

Arrow's impossibility theorem. Definitions

Definition

An aggregation rule $f : (\text{Ord}(A))^n \rightarrow \text{Ord}(A)$ satisfies

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An aggregation rule $f : (\text{Ord}(A))^n \rightarrow \text{Ord}(A)$ satisfies

U (unanimity) iff

$$(\forall a, b \in A) ((\forall i < n) a \prec_i b) \rightarrow a f(\pi) b$$

for any profile $\pi = (\prec_0, \prec_1, \dots, \prec_{n-1})$ in $(\text{Ord}(A))^n$;

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for any profile $\pi = (\prec_0, \prec_1, \dots, \prec_{n-1})$ in $(\text{Ord}(A))^n$;

IIA (independence of irrelevant alternatives) iff

$$(\forall a, b \in A) ((\forall i < n) a \prec_i b \leftrightarrow a \prec'_i b) \rightarrow (a f(\pi) b \leftrightarrow a f(\pi') b)$$

for all profiles $\pi = (\prec_0, \prec_1, \dots, \prec_{n-1})$, $\pi' = (\prec'_0, \prec'_1, \dots, \prec'_{n-1})$ in $(\text{Ord}(A))^n$.

Arrow's impossibility theorem

An aggregation rule $f : (\text{Ord}(A))^n \rightarrow \text{Ord}(A)$ satisfies

D (dictatorship) iff f is a projection, i.e. there is $i < n$ such that for all $\pi = (\prec_0, \prec_1, \dots, \prec_{n-1}) \in (\text{Ord}(A))^n$

$$f(\pi) = \prec_i .$$

Theorem (K. Arrow, 1950,1963)

For any natural number $n \geq 1$, finite set A of cardinality $|A| \geq 3$, and aggregation rule $f : (\text{Ord}(A))^n \rightarrow \text{Ord}(A)$ if f satisfies **U** and **IIA** then f satisfies **D**.

Arrow's impossibility theorem in terms of choice functions. Notation and definitions

- A – a (finite) set (of alternatives);
- $[A]^2 = \{B \subseteq A : |B| = 2\}$;
- 2-choice function on A – function $c : [A]^2 \rightarrow A$ satisfying

$$(\forall p \in [A]^2) c(p) \in p.$$

Definition

A 2-choice function c on A is *rational* iff there is a (strict) linear order \prec on A such that $c(p) = \max_{\prec} p$ for all $p \in [A]^2$, i.e.

$$c(\{a, b\}) = b \leftrightarrow a \prec b$$

for all different $a, b \in A$.

Arrow's impossibility theorem in terms of choice functions. Notation and definitions

- *(individual) preferences* = rational 2-choice function on A ;
- $\mathfrak{R}_2(A)$ – the set of all rational 2-choice functions on A ;
- n – a natural number (of voters);
- *profile* = n -tuple of rational 2-choice functions on A ;
- *(universal) aggregation rule* = function $f : (\mathfrak{R}_2(A))^n \rightarrow \mathfrak{R}_2(A)$.

Note

The function Θ that assigns to each strict linear order \prec the rational 2-choice function \max_{\prec} is a bijection between $\text{Ord}(A)$ and $\mathfrak{R}_2(A)$.

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U (unanimity) iff

$$(\forall p \in [A]^2) (\forall b \in p) ((\forall i < n) \mathbf{c}(p) = b) \rightarrow f(\pi)(p) = b$$

for any profile $\pi = (\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})$ in $(\mathfrak{R}_2(A))^n$;

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for any profile $\pi = (\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})$ in $(\mathfrak{R}_2(A))^n$;

IIA (independence of irrelevant alternatives) iff

$$(\forall p \in [A]^2) ((\forall i < n) \mathbf{c}_i(p) = \mathbf{c}'_i(p)) \rightarrow f(\pi)(p) = f(\pi')(p)$$

for all profiles $\pi = (\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})$, $\pi' = (\mathbf{c}'_0, \mathbf{c}'_1, \dots, \mathbf{c}'_{n-1})$ in $(\mathfrak{R}_2(A))^n$.

Arrow's impossibility theorem in terms of choice functions. Notation and definitions

Note

A function $f : (\mathfrak{R}_2(A))^n \rightarrow \mathfrak{R}_2(A)$ satisfies **IIA** iff for all $p \in [A]^2$ there is a function $f_p : p^n \rightarrow p$ such that

$$f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(p) = f_p(\mathbf{c}_0(p), \mathbf{c}_1(p), \dots, \mathbf{c}_{n-1}(p))$$

for all $p \in [A]^2$ and $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{R}_2(A)$.

A function $f : (\mathfrak{R}_2(A))^n \rightarrow \mathfrak{R}_2(A)$ satisfies **IIA** and **U** iff there is a conservative function $\hat{f} : A^n \rightarrow A$ such that

$$f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(p) = \hat{f}(\mathbf{c}_0(p), \mathbf{c}_1(p), \dots, \mathbf{c}_{n-1}(p)).$$

for all $p \in [A]^2$ and $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{R}_2(A)$.

Conservative functions

Definition

A function $g : A^n \rightarrow A$ is *conservative* iff

$$(\forall x_0, x_1, \dots, x_{n-1} \in A) \bigvee_{i < n} (g(x_0, x_1, \dots, x_{n-1}) = x_i).$$

Arrow's impossibility theorem in terms of choice functions

An aggregation rule $f : (\text{Ord}(A))^n \rightarrow \text{Ord}(A)$ satisfies

D iff f is a projection, i.e. there is $i < n$ such that for all $\pi = (\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1}) \in (\mathfrak{R}_2(A))^n$

$$f(\pi) = \mathbf{c}_i.$$

Theorem

For any natural number $n \geq 1$, finite set A of cardinality $|A| \geq 3$, and aggregation rule $f : (\mathfrak{R}_2(A))^n \rightarrow \mathfrak{R}_2(A)$ if f satisfies **U** and **IIA** then f satisfies **D**.

Shelah's extension

Shelah's extension. Notation and definitions

- A – a non-empty (finite) set (of alternatives);
- r – a natural number (technical parameter), $r \geq 1$;
- $[A]^r = \{B \subseteq A : |B| = r\}$;
- r -choice function on A – function $\mathfrak{c} : [A]^r \rightarrow A$ satisfying

$$(\forall p \in [A]^r) \mathfrak{c}(p) \in p.$$
- $\mathfrak{C}_r(A)$ – the set of all r -choice function on A ;

Definition

A set $\mathfrak{D} \subseteq \mathfrak{C}_r(A)$ is *symmetric* if for any function $\mathfrak{c} \in \mathfrak{D}$ and permutation $\sigma \in S_A$ the function \mathfrak{c}_σ defined by

$$(\forall p \in [A]^r) \mathfrak{c}_\sigma(p) = \sigma^{-1}\mathfrak{c}(\sigma p),$$

belongs to \mathfrak{D} .

Symmetric sets of r -choice functions

Exemples

- The set $\mathfrak{R}_2(A)$.
- The set of all function $\mathfrak{c} \in \mathfrak{C}_r(A)$ such that $\mathfrak{c}(p)$ is the median element in p according to some ordering (r is odd).
- The set $\{\mathfrak{c} \in \mathfrak{C}_2(A) : (\exists x \in A)(\forall y \in A \setminus \{x\}) \mathfrak{c}(\{x, y\}) = x\}$.
- Let \prec be a strict partial order on A and $\mathfrak{C}_r^\prec(A)$ a set of all functions $\mathfrak{c} \in \mathfrak{C}_r(A)$ such that $\mathfrak{c}(p)$ is some non-dominated element of p , i.e.

$$(\forall x \in p) \mathfrak{c}(p) \not\prec x.$$

Let W be a set of strict partial order on A closed under isomorphisms. The set $\bigcup_{\prec \in W} \mathfrak{C}_r^\prec(A)$ is symmetric.

Shelah's extension. Notation and definitions

- (*individual*) preferences = r -choice function on A ;
- n – a natural number (of voters), $n \geq 1$;
- *profile* = n -tuple of r -choice functions on A ;
- *aggregation rule* = function $f : (\mathfrak{C}_r(A))^n \rightarrow \mathfrak{C}_r(A)$;
- $\mathcal{V}(A, r)$ – the set of all aggregation rules (of all arity $n \geq 1$).

Shelah's extension. Notation and definitions

Definition

An n -ary aggregation rule $f \in \mathcal{V}(A, r)$ is *normal* iff

- (i) $f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(q) \in \{\mathbf{c}_0(q), \mathbf{c}_1(q), \dots, \mathbf{c}_{n-1}(q)\}$
for all $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{C}_r(A)$ and $q \in [A]^r$;

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for all $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{C}_r(A)$ and $q \in [A]^r$;
- (ii) $(\mathbf{c}_0(q), \mathbf{c}_1(q), \dots, \mathbf{c}_{n-1}(q)) = (\mathbf{c}'_0(q), \mathbf{c}'_1(q), \dots, \mathbf{c}'_{n-1}(q)) \rightarrow$
 $f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(q) = f(\mathbf{c}'_0, \mathbf{c}'_1, \dots, \mathbf{c}'_{n-1})(q)$
for all $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1}, \mathbf{c}'_0, \mathbf{c}'_1, \dots, \mathbf{c}'_{n-1} \in \mathfrak{C}_r(A)$ and $q \in [A]^r$.

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for all $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{C}_r(A)$ and $q \in [A]^r$;
- (ii) $(\mathbf{c}_0(q), \mathbf{c}_1(q), \dots, \mathbf{c}_{n-1}(q)) = (\mathbf{c}'_0(q), \mathbf{c}'_1(q), \dots, \mathbf{c}'_{n-1}(q)) \rightarrow$
 $f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(q) = f(\mathbf{c}'_0, \mathbf{c}'_1, \dots, \mathbf{c}'_{n-1})(q)$
for all $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1}, \mathbf{c}'_0, \mathbf{c}'_1, \dots, \mathbf{c}'_{n-1} \in \mathfrak{C}_r(A)$ and $q \in [A]^r$.

- $\mathcal{N}(A, r)$ – the set of all normal aggregation rules in $\mathcal{V}(A, r)$.

Shelah's extension. Notation and definitions

Definition

An n -ary aggregation rule $f \in \mathcal{V}(A, r)$ is *simple* iff

$$\begin{aligned}
 (\mathbf{c}_0(p), \mathbf{c}_1(p), \dots, \mathbf{c}_{n-1}(p)) = (\mathbf{c}_0(q), \mathbf{c}_1(q), \dots, \mathbf{c}_{n-1}(q)) &\rightarrow \\
 f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(p) = f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(q) & \\
 \text{for all } \mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{C}_r(A) \text{ and } p, q \in [A]^r. &
 \end{aligned}$$

Shelah's extension. Notation and definitions

Definition

An n -ary aggregation rule $f \in \mathcal{V}(A, r)$ is *simple* iff

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 (\mathbf{c}_0(p), \mathbf{c}_1(p), \dots, \mathbf{c}_{n-1}(p)) = (\mathbf{c}_0(q), \mathbf{c}_1(q), \dots, \mathbf{c}_{n-1}(q)) &\rightarrow \\
 f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(p) = f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(q) & \\
 \text{for all } \mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{C}_r(A) \text{ and } p, q \in [A]^r. &
 \end{aligned}$$

- $\mathcal{S}(A, r)$ – the set of all simple aggregation rules in $\mathcal{V}(A, r)$.

Shelah's extension. Notation and definitions

Definition

An n -ary aggregation rule $f \in \mathcal{V}(A, r)$ is a *dictatorship* iff it is a projection, i.e. iff there is $i < n$ such that $f(\mathfrak{c}_0, \mathfrak{c}_1, \dots, \mathfrak{c}_{n-1}) = \mathfrak{c}_i$ for all $\mathfrak{c}_0, \mathfrak{c}_1, \dots, \mathfrak{c}_{n-1} \in \mathfrak{C}_r(A)$

Shelah's extension. Notation and definitions

Definition

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- $\mathcal{E}(A, r)$ – the set of all dictatorships in $\mathcal{V}(A, r)$.

Shelah's extension. Notation and definitions

Note

An n -ary function $f \in \mathcal{V}(A, r)$ is normal iff for all $p \in [A]^2$ there is a conservative function $f_p : p^n \rightarrow p$ such that

$$f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(p) = f_p(\mathbf{c}_0(p), \mathbf{c}_1(p), \dots, \mathbf{c}_{n-1}(p))$$

for all $p \in [A]^2$ and $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{C}_r(A)$.

An n -ary function $f : (\mathfrak{R}_2(A))^n \rightarrow \mathfrak{R}_2(A)$ is normal and simple iff there is a conservative function $\widehat{f} : A^n \rightarrow A$ such that

$$f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(p) = \widehat{f}(\mathbf{c}_0(p), \mathbf{c}_1(p), \dots, \mathbf{c}_{n-1}(p)).$$

for all $p \in [A]^2$ and $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{C}_r(A)$.

Shelah's extension. Notation and definitions

Definition

An n -ary aggregation rule $f \in \mathcal{V}(A, r)$ *preserves* a set $\mathcal{D} \subseteq \mathfrak{C}_r(A)$ (or f is a *polymorphism* of \mathcal{D}) and \mathcal{D} is *preserved* under f iff

$$f(\mathbf{c}_1, \mathbf{c}_2, \dots, \mathbf{c}_n) \in \mathcal{D} \text{ for all } \mathbf{c}_1, \mathbf{c}_2, \dots, \mathbf{c}_n \in \mathcal{D}.$$

The set of all $f \in \mathcal{V}(A, r)$ that preserves $\mathcal{D} \subseteq \mathfrak{C}_r(A)$ is denoted by $\text{pol } \mathcal{D}$.

Shelah's extension. Notation and definitions

Definition

A set $\mathfrak{D} \subseteq \mathfrak{C}_r(A)$

- *has the Arrow property iff*

$$\text{pol } \mathfrak{D} \cap \mathcal{N}(A, r) = \mathcal{E}(A, r).$$

- *has the simple Arrow property iff*

$$\text{pol } \mathfrak{D} \cap \mathcal{N}(A, r) \cap \mathcal{S}(A, r) = \mathcal{E}(A, r).$$

Shelah's theorem on the Arrow property

Theorem (S. Shelah, 2005)

Let A be a finite set. Then there are natural numbers r_1, r_2 (e.g. $r_1 = r_2 = 7$) such that for any natural number r , $r_1 \leq r \leq |A| - r_2$, any non-empty proper symmetric subset \mathcal{D} of the set $\mathfrak{C}_r(A)$ has the Arrow property.

Complete classification of symmetric sets of r -choice function without the Arrow property

Exceptional cases: $\mathfrak{C}_3^K(A)$

Let $|A| = 4$ and let K be the *Klein four-group* of permutations of A . For any sets $p, q \in [A]^3$ there is only one permutation $\sigma_{p,q} \in K$ for which

$$q = \sigma_{p,q}(p).$$

- $\mathfrak{C}_3^K(A)$ is the set of all functions $\mathfrak{c} \in \mathfrak{C}_3(A)$ such that

$$\mathfrak{c}(q) = \sigma_{p,q}\mathfrak{c}(p) \text{ for all } p, q \in [A]^3.$$

Exceptional cases: $\mathfrak{C}_3^K(A)$

The set $\mathfrak{C}_3^K(A)$ is symmetric and it contains exactly three elements $\mathfrak{c}_0, \mathfrak{c}_1, \mathfrak{c}_2$ (we denote $A = \{a, b, c, d\}$):

q	$\mathfrak{c}_0(q)$	$\mathfrak{c}_1(q)$	$\mathfrak{c}_2(q)$
$\{a, b, c\}$	a	b	c
$\{a, b, d\}$	b	a	d
$\{a, c, d\}$	c	d	a
$\{b, c, d\}$	d	c	b

Exceptional cases: $\mathfrak{C}_3^K(A)$

Let the conservative function $w: A^2 \rightarrow A$ be defined by

w	a	b	c	d
a	a	a	c	d
b	b	b	c	d
c	a	b	c	c
d	a	b	d	d

Let the 2-ary simple normal aggregation rule $f \in \mathcal{V}(A, 3)$ be defined by

$$f_q(x, y) = w(x, y)$$

for all $q \in [A]^3$ and $x, y \in q$.

Exceptional cases: $\mathfrak{C}_3^K(A)$

We have

f	c_0	c_1	c_2
c_0	c_0	c_0	c_2
c_1	c_1	c_1	c_2
c_2	c_0	c_1	c_2

Proposition

The set $\mathfrak{C}_3^K(A)$ is preserved by the normal and simple aggregation rule f . So, $\mathfrak{C}_3^K(A)$ is non-empty proper symmetric subset of the set $\mathfrak{C}_3(A)$ without the simple Arrow property.

Exceptional cases: $\mathfrak{C}_2^i(A)$

Let $r = 2$ and $|A| \geq 2$. For any $a \in A$, $i \in \{0, 1\}$ and $\mathbf{c} \in \mathfrak{C}_2(A)$

$$Z_a^{\mathbf{c}} = \{b \in A \setminus \{a\} : \mathbf{c}(\{a, b\}) = a\},$$

$$W_i^{\mathbf{c}} = \{a \in A : |Z_a^{\mathbf{c}}| = i \pmod{2}\}.$$

- $\mathfrak{C}_2^i(A) = \{\mathbf{c} \in \mathfrak{C}_2(A) : W_{(1-i)}^{\mathbf{c}} = \emptyset\}.$

Example

$$A = \{a, b, c\}, \quad \mathfrak{C}_2^1(A) = \{\mathbf{c}_0, \mathbf{c}_1\}.$$

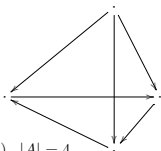
q	$\mathbf{c}_0(q)$	$\mathbf{c}_1(q)$
$\{a, b\}$	b	a
$\{b, c\}$	c	b
$\{a, c\}$	a	c

Exceptional cases: $\mathfrak{C}_2^i(A)$

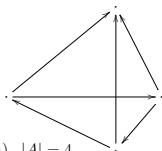
Another definition. Any function $\mathfrak{c} \in \mathfrak{C}_2(A)$ may be represented by the *tournament* $\Gamma_{\mathfrak{c}} = (A, E)$ where

$$E = \{(a, b) \in A^2 : a \neq b \wedge \mathfrak{c}(\{a, b\}) = b\}.$$

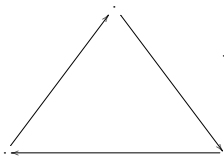
- The sets $\mathfrak{C}_2^0(A)$ and $\mathfrak{C}_2^1(A)$ are the sets of all functions $\mathfrak{c} \in \mathfrak{C}_2(A)$ such that the *indegree* of any node of the tournament $\Gamma_{\mathfrak{c}}$ is even (respectively, odd).



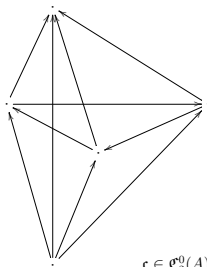
$c \in \mathfrak{C}_2^0(A)$, $|A| = 4$



$c \in \mathfrak{C}_2^1(A)$, $|A| = 4$



$c \in \mathfrak{C}_2^1(A)$, $|A| = 3$



$c \in \mathfrak{C}_2^0(A)$, $|A| = 5$

Exceptional cases: $\mathfrak{C}_2^i(A)$

For any non-empty set A the 3-ary normal simple aggregation rule $\ell^A \in \mathcal{V}(A, 2)$ is defined by

$$\ell_q^A(x, x, y) = \ell_q^A(x, y, x) = \ell_q^A(y, x, x) = y$$

for all $q \in [A]^2$ and $x, y \in q$.

Proposition

Each of the sets $\mathfrak{C}_2^0(A)$, $\mathfrak{C}_2^1(A)$, and $\mathfrak{C}_2^0(A) \cup \mathfrak{C}_2^1(A)$ is preserved by the normal and simple aggregation rule ℓ^A . So, it does not have the simple Arrow property. Besides, each of these sets is symmetric, and

- 1 $\mathfrak{C}_2^0(A) \neq \emptyset \Leftrightarrow |A| = 0, 1 \pmod{4}$;
- 2 $\mathfrak{C}_2^1(A) \neq \emptyset \Leftrightarrow |A| = 0, 3 \pmod{4}$;
- 3 $\mathfrak{C}_2^0(A) \cup \mathfrak{C}_2^1(A) \neq \mathfrak{C}_2(A)$.

Complete classification of symmetric sets of r -choice function without the Arrow property

Theorem (N. Polyakov, 2014)

Let A be a finite set, r be a natural number, and \mathfrak{D} be a non-empty proper symmetric subset of the set $\mathfrak{C}_r(A)$. Then the set \mathfrak{D} does not has the Arrow property if and only if one of the following conditions holds:

- ① $r = 2$, $|A|$ equals 0 or 1 (mod 4), and $\mathfrak{D} = \mathfrak{C}_2^0(A)$,
- ② $r = 2$, $|A|$ equals 0 or 3 (mod 4), and $\mathfrak{D} = \mathfrak{C}_2^1(A)$,
- ③ $r = 2$, $|A| = 0$ (mod 4), and $\mathfrak{D} = \mathfrak{C}_2^0(A) \cup \mathfrak{C}_2^1(A)$,
- ④ $r = 3$, $|A| = 4$, and $\mathfrak{D} = \mathfrak{C}_3^K(A)$.

Arrow property for classes of decision rules

Arrow property for classes of decision rules.

Notation and definitions

- Q – a non-empty (finite) set (of conditions);
- A – a non-empty (finite) set (of solutions);
- $\mathfrak{D} \subseteq {}^Q A$ – a set (of decision rules).

Definition

A set $\mathfrak{D} \subseteq {}^Q A$ is (weakly) *symmetric* iff for any permutation σ of A there is a permutation σ^* of Q such that for all function $\mathfrak{c} \in \mathfrak{D}$ the functions \mathfrak{c}_σ and $\tilde{\mathfrak{c}}_\sigma$ defined by

$$\mathfrak{c}_\sigma(q) = \sigma^{-1} \mathfrak{c}(\sigma^* q) \quad \text{and} \quad \tilde{\mathfrak{c}}_\sigma(q) = \sigma \mathfrak{c}((\sigma^*)^{-1} q)$$

belong to \mathfrak{D} .

Arrow property for classes of decision rules.

Notation and definitions

Examples

- $Q = [A]^r$, $\mathfrak{D} = \mathfrak{C}_r(A)$;
- $Q = \mathcal{P}(A) \setminus \{\emptyset\}$, \mathfrak{D} is a set of all (total) choice functions;
- Q is a set of subsets of A enriched with some additional structure, for example
 - Q is a set of all linear orders on A ,
 - Q is a set of non-empty multisets such that the underlying set of elements is a subset of A ,
 and \mathfrak{D} is a set of choice function;
- Q is a set of all linear orders on some set B , $A = \mathcal{P}(B)$, $\mathfrak{D} = {}^Q A$.

Arrow property for classes of decision rules.

Notation and definitions

- n – a natural number (of voters), $n \geq 1$;
- *profile* = n -tuple of decision rules in \mathfrak{D} ;
- (*simple*) *aggregation rule* = conservative function $f : A^n \rightarrow A$;
- $\mathcal{V}(A)$ – the set of all simple aggregation rules (of all arity $n \geq 1$).

Arrow property for classes of decision rules.

Notation and definitions

Definition

A function $f : A^n \rightarrow A$ *preserves* a set $\mathfrak{D} \subseteq {}^Q A$ and \mathfrak{D} is *preserved* under f iff for all $\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{D}$ the set \mathfrak{D} contains the function $f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})$ defined by

$$f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1})(q) = f(\mathbf{c}_0(q), \mathbf{c}_1(q), \dots, \mathbf{c}_{n-1}(q))$$

for all $q \in Q$.

Arrow property for classes of decision rules.

Notation and definitions

- $\text{pol } \mathfrak{D}$ – the set of all functions $f : A^n \rightarrow A$ (of any arity n) that preserve a set $\mathfrak{D} \subseteq {}^Q A$;
- $\text{inv}_Q f$ – the set of all sets $\mathfrak{D} \subseteq {}^Q A$ that is preserved under a function $f : A^n \rightarrow A$;
- for all sets $\mathcal{F} \subseteq \bigcup_{n < \omega} A^n A$ and $\mathbb{D} \subseteq \mathcal{P}({}^Q A)$

$$\text{inv}_Q \mathcal{F} = \bigcap_{f \in \mathcal{F}} \text{inv}_Q f \quad \text{and} \quad \text{pol } \mathbb{D} = \bigcap_{\mathfrak{D} \in \mathbb{D}} \text{pol } \mathfrak{D}.$$

Arrow property for classes of decision rules.

Notation and definitions

Theorem

The couple $(\text{inv}_Q, \text{pol})$ is a Galois correspondence between the Boolean lattices $\mathcal{P}(\bigcup_{n < \omega} A^n A)$ and $\mathcal{P}(\mathcal{P}(Q A))$. Galois-closed sets $\mathcal{F} \subseteq \bigcup_{n < \omega} A^n A$ are closed under composition and contain all projections, i.e. is clones. If a set $\mathcal{D} \subseteq Q A$ is symmetric, then the clone $\text{pol } \mathcal{D}$ is symmetric, i.e. for all permutation σ of A and n -ary function $f \in \text{pol } \mathcal{D}$ the clone $\text{pol } \mathcal{D}$ contains a function $f_\sigma : A^n \rightarrow A$, defined by

$$f_\sigma(\mathbf{a}) = \sigma^{-1} f(\sigma \mathbf{a})$$

for all $\mathbf{a} \in A^n$.

Arrow property for classes of decision rules.

Notation and definitions

Definition

The set $\mathfrak{D} \subseteq {}^Q A$ has a (simple) *Arrow property* iff for any natural number n and any n -ary function $f \in \text{pol } \mathfrak{D} \cap \mathcal{V}(A)$ there is a number i ($i < n$) for which

$$(\forall \mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1} \in \mathfrak{D}) f(\mathbf{c}_0, \mathbf{c}_1, \dots, \mathbf{c}_{n-1}) = \mathbf{c}_i,$$

i.e.

$$f(\mathbf{a}) = a_i,$$

for all n -tuples $\mathbf{a} \in \{(\mathbf{c}_0(q), \mathbf{c}_1(q), \dots, \mathbf{c}_{n-1}(q)) : q \in Q, \mathbf{c}_i \in \mathfrak{D}\}$.

Are there *many* symmetric sets of decision rules without Arrow properties?

Example

Let A be a finite set, r be a natural number, $2 \leq r \leq |A| - 1$, and \mathfrak{D} be a symmetric subset of $\mathfrak{C}_r(A)$. Let

- $Q = [A]^r \cup [A]^{r+1}$;
- \mathfrak{C} be the set of all choice functions c on Q such that $c \upharpoonright [A]^r \in \mathfrak{D}$.

Are there *many* symmetric sets of decision rules without Arrow properties?

Example

Then

- \mathfrak{C} is symmetric, $\mathfrak{C} \upharpoonright [A]^r = \mathfrak{D}$;
- \mathfrak{C} is preserved under the conservative function $f : A^{r+1} \rightarrow A$ defined by

$$f(\mathbf{x}) = \begin{cases} x_1, & \text{if } |\text{ran } \mathbf{x}| = r + 1; \\ x_0, & \text{if } |\text{ran } \mathbf{x}| \leq r. \end{cases}$$

for all $\mathbf{x} = (x_0, x_1, \dots, x_r) \in A^{r+1}$.

Are there *many* symmetric sets of decision rules without Arrow properties?

Example

q	$c_0(q)$	$c_1(q)$	\dots	$c_r(q)$	$f(c_0, c_1, \dots, c_r)(q)$
$\{a_0, a_1, \dots, a_r\}$	a_0	a_1	\dots	a_r	a_1
$\{a_0, a_1, \dots, a_{r-1}\}$	a_0	a_1	\dots	a_1	a_0

\mathfrak{C} does not have the Arrow property.

Are there *many* symmetric sets of decision rules without Arrow properties?

Theorem

Let A and Q be a finite sets, and \mathfrak{D} be a symmetric subset of ${}^Q A$. Then there are finite sets $B \supseteq A$ and $P \supseteq Q$ and a symmetric set $\mathfrak{C} \subseteq {}^P B$ such that

- \mathfrak{C} does not have the Arrow property;
- ${}^Q A \cap \mathfrak{C} \upharpoonright Q = \mathfrak{D}$.

Arrow property for classes of decision rules.

Notation and definitions

- \mathbb{B}_0 is the set of all sets $\mathfrak{B} \subseteq {}^Q A$ of the form

$$\{\mathbf{c} \in {}^Q A : \mathbf{c}(q) \in B\}$$

where $q \in Q$ and $B \subseteq A$;

- \mathbb{B}_1 is the set of all sets $\mathfrak{B} \subseteq {}^Q A$ of the form

$$\{\mathbf{c} \in {}^Q A : \mathbf{c}(p) = a \vee \mathbf{c}(q) = b\},$$

where $p, q \in Q$ and $a, b \in A$;

Arrow property for classes of decision rules.

Notation and definitions

- $\mathbb{B}_2(R)$ is the set of all sets $\mathfrak{B} \subseteq {}^Q A$ of the form

$$\{\mathbf{c} \in {}^Q A : \mathbf{c}(q) = \sigma \mathbf{c}(p)\},$$

where R is a binary relation on $A^{<\omega}$, $p, q \in Q$, $\sigma \in S_A$ and $(\mathbf{b}, \sigma \mathbf{b}) \in R$ for all $\mathbf{b} \in A^{<\omega}$;

- $\mathbb{B}_3(\Pi)$ is the set of all sets $\mathfrak{B} \subseteq {}^Q A$ of the form

$$\{\mathbf{c} \in {}^Q A : \mathbf{c} \upharpoonright P \in \text{inv}_P(\Pi_B)\},$$

where Π is a Post's class closed under duality, $B \in [A]^2$, Π_B is a clone on B naturally isomorphic to Π , $P \subseteq Q$.

Arrow property for classes of decision rules.

Notation and definitions

Note

There are only six Post's classes $\Pi \subseteq T_{01}$ closed under duality:
 $O_1, D_1, D_2, L_4, A_4, T_{01}$.

Definition

A binary relation R on $A^{<\omega}$ is *stable* iff

- ① $\mathbf{a} R \mathbf{b} \rightarrow \mathbf{a} = \sigma \mathbf{b}$ for some permutation σ of A ;
- ② $\mathbf{a} R \mathbf{b} \rightarrow \sigma \mathbf{a} \tau R \sigma \mathbf{b} \tau$ for any permutation σ of A , natural number k and function $\tau: \{0, 1, \dots, k-1\} \rightarrow \text{dom } \mathbf{a}$.

Arrow property for classes of decision rules.

Notation and definitions

Let $\mathfrak{D} \subseteq {}^Q A$. For any natural number r

- $r(\mathfrak{D}) = \max_{q \in Q} |\{\mathbf{c}(q) : \mathbf{c} \in \mathfrak{D}\}|$
- $Q_{\mathfrak{D},r} = \{q \in Q : |\{\mathbf{c}(q) : \mathbf{c} \in \mathfrak{D}\}| \leq r\};$
- $\mathfrak{D}_r = \mathfrak{D} \upharpoonright Q_{\mathfrak{D},r}$
- $\mathfrak{D}_r^+ = \{\mathbf{c} \in {}^Q A : \mathbf{c} \upharpoonright Q_{\mathfrak{D},r} \in \mathfrak{D}_r\};$

Arrow property for classes of decision rules

Theorem

Let A and Q be non-empty finite sets. Let $\mathfrak{D} \subseteq {}^Q A$ be a symmetric set without Arrow property. Then there are a stable binary relation R on $A^{<\omega}$, a Post's class $\Pi \in \{O_1, D_1, D_2, L_4, A_4, T_{01}\}$ and a set $\mathbb{B} \subseteq \mathbb{B}_0 \cup \mathbb{B}_1 \cup \mathbb{B}_2(R) \cup \mathbb{B}_3(\Pi)$ such that one of two following conditions holds

- 1 $\mathfrak{D} = \bigcap \mathbb{B}$;
- 2 there is a natural number $r < r(\mathfrak{D})$ such that $\mathfrak{D} = \mathfrak{D}_r^+ \cap \bigcap \mathbb{B}$, and any n -ary function $f \in \text{pol } \mathfrak{D} \cap \mathcal{V}(A)$ coincides with a projection on the set $A_{\leq r}^n = \{\mathbf{a} \in A^n : |\text{ran } \mathbf{a}| \leq r\}$ (hence \mathfrak{D}_r has the Arrow property).

Some positive results

Theorem

Let A be a non-empty finite set, $|A| \geq 3$, and f be an arbitrary non-dictatorship function in the clone \mathcal{D} generated by a (majority) function ∂ satisfying

$$\partial(x, x, y) = \partial(x, y, x) = \partial(y, x, x) = x.$$

Then a set $\mathfrak{D} \subseteq \mathfrak{C}_2(A)$ belong of $\text{inv}_{[A]^2} f$ if and only if the set \mathfrak{D} is an intersection of a family of sets of the form

$$\{\mathbf{c} \in \mathfrak{C}_2(A) : \mathbf{c}(\{a, b\}) = a \rightarrow \mathbf{c}(\{c, d\}) = d\}.$$

where $a, b, c, d \in A$.

Thank you!